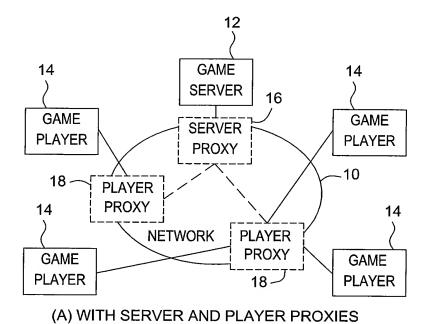
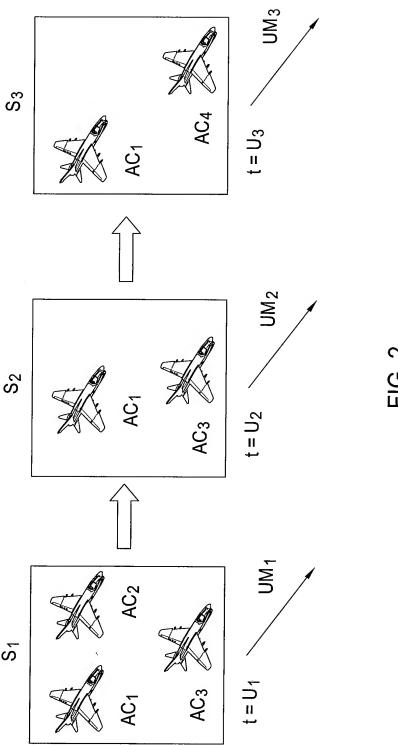


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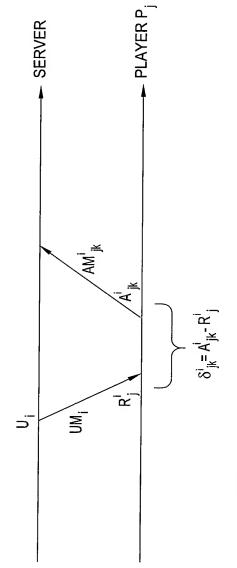


DISTRIBUTED GAME ENVIRONMENT

FIG. 1



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MESSAGE EXCHANGE BETWEEN SERVER AND PLAYERS

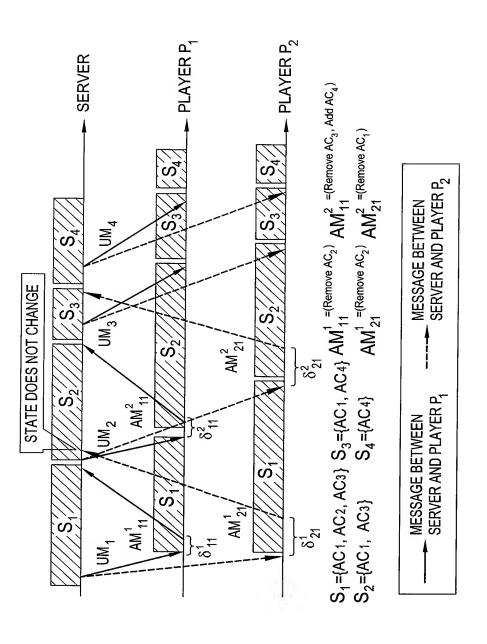
ALGORITHM FAIR-ORDER MESSAGE QUEUEING (ACTION_MESSAGE M_k): 1: COMPUTE D(M_k) = DELIVERY TIME OF M_k ;

2: INSERT M_k INTO DELIVERY QUEUE SORTED ACCORDING TO $D(M_k)$; 3: IF (DELIVERY QUEUE SIZE > 1)

RECOMPUTE DELIVERY TIME OF EXISTING MESSAGES;

ALGORITHM FOR FAIR-ORDER MESSAGE QUEUEING

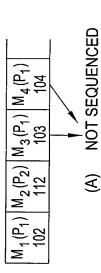
FIG. 5



FAIR-ORDER MESSAGE DELIVERY FOR STATE TRANSITIONS SHOWN IN FIG. 2

FIG. 4

MESSAGE FROM P₁ WITH SEQ. NO 101 HAS NOT ARRIVED.



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MESSAGE FROM P₁ WITH SEQ. NO 101 HAS NOT ARRIVED BUT DELIVERY TIME OF M₁ IS REACHED AND M₁ IS DELIVERED.

 $M_2(P_2) M_3(P_1) M_4(P_1)$ 112 103 104 (B) SEQUENCED

EXAMPLE WHERE MESSAGES ARRIVE AFTER THEIR WAIT TIMEOUT.

FIG. 6

ALGORITHM FAIR-ORDER MESSAGE DEQUEUING (ACTION_MESSAGE M_k): 1: DELIVERY M_k AT $D(M_k$);

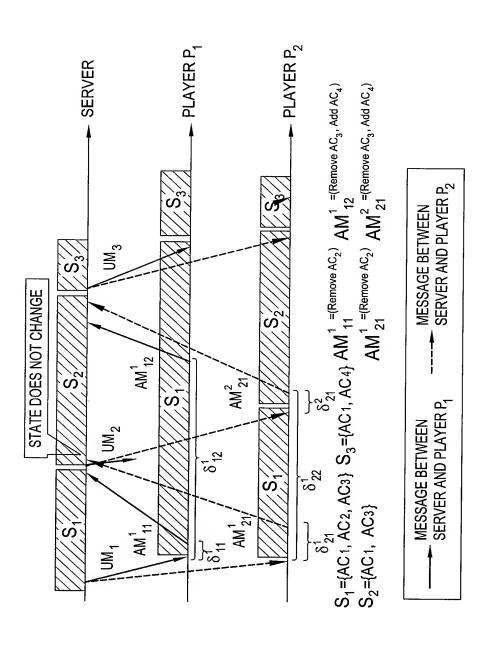
2: IF (DELIVERY QUEUE SIZE > 1)
3: RECOMPUTE DELIVERY TI

RECOMPUTE DELIVERY TIME OF EXISTING MESSAGES;

ADDITIONAL ALGORITHM FOR FAIR-ORDER MESSAGE DEQUEUING WHEN MESSAGES DO NOT ARRIVE WITHIN THEIR WAIT TIMEOUT.

-<u>|</u>G. 7

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EXAMPLE OF AN INCONSISTENT VIEW OF THE GAME BETWEEN TWO PLAYERS, WHERE THE SEQUENCE OF STATE CHANGES AT THE SERVER IS SHOWN IN FIG. 2

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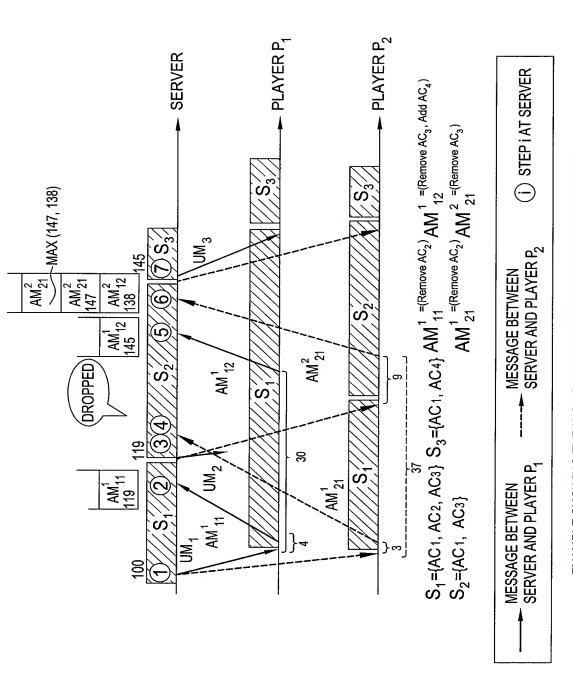
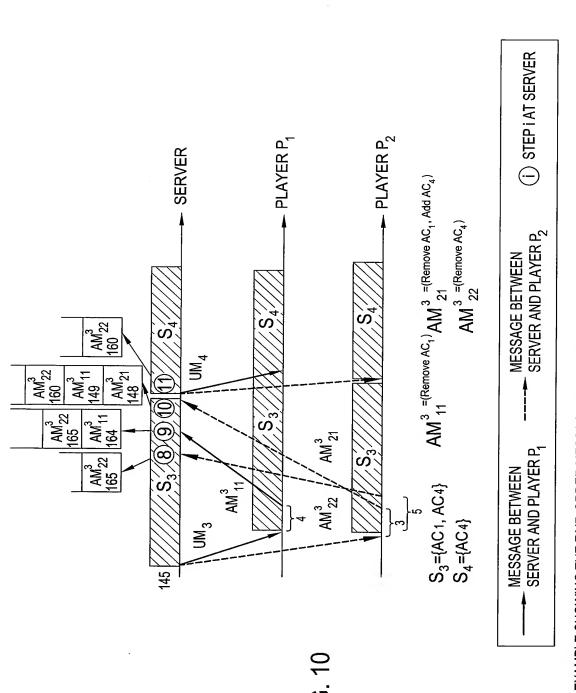


FIG. 9

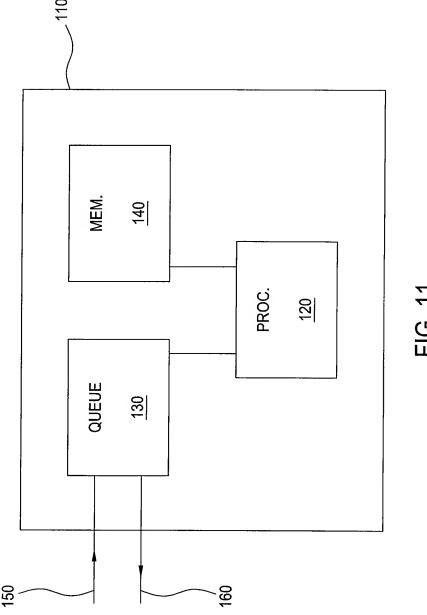
EXAMPLE SHOWING THE FAIR-ORDER MESSAGE DELIVERY ALGORITHM.

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EXAMPLE SHOWING THE FAIR-ORDER MESSAGE DELIVERY ALGORITHM WITH OUT-OF-ORDER MESSAGE RECEPTION.

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